

Android Security and the Elusive HSM

Mobile Digital Wallet Security Summit

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Android Security and the Elusive HSM – Jon Oberheide

Slide #1

Slide #2

Introduction

- Jon Oberheide
- CTO, Duo Security
- Today
 - High level look at Android security and how HSMs can help
 - Only 30 minutes!
 Lots of external refs!







Use case



- Mobile platform security is important (surprise!)
- Apps have data to protect (confidentiality/integrity)
- Examples
 - Duo Mobile 2FA app
 - Visa Mobile app? V.me?







- Attacker wants access to sensitive data
 - Stages of attack/capability
 - Points of attack disruption/mitigation







- Introduction
- Gaining code execution
- Escalating privileges
- Recent HSM developments
- Wrap-up





- First step is getting a foothold on the mobile device
 - Need code execution on victim's device
- A couple easy vectors for attackers:
 - Social vector: malicious applications
 - Technical vector: exploitation of existing applications

PoC malicious apps



In the past, mostly researcher PoCs











• Nowadays, real-world malware is out there

Security ... A NBCNEWS.com

New DroidDream malware infects 24 Android malware numbers

DroidDreamLight Trojan has a

explode to 25,000 in June 2012

IT administrators await mammoth M…

Rapper Soulja Boy blames Facebo... ▶ malware threats found has hit 25,000. In Ju

malware threats found has hit 25,000. In Ju 10,000, easily the largest find for a month y

Plankton malware drifts into Android Market

Join thousands of others, and sign up for

VISA

FAST FEED

Android Malware Infects 500,000 Users

Android S BY NEAL UNGERLEIDER | AUGUST 20, 2012

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Stopping malicious apps

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- Mobile antivirus
 - Reactive signature-based detection of malicious apps



• Mobile malware exists, but these folks tend to push the hype a bit beyond reality





- Google's Bouncer
 - Guards the entry point to the Android Market
 - Dynamically analyze submitted apps and block malicious apps from being published
- Dynamic analysis is <u>hard</u>
- Fingerprinting Bouncer is <u>easy</u>
- Catches some malware, but easy to bypass

http://jon.oberheide.org/files/summercon12-bouncer.pdf









- Threat of client-side applications
 - Large attack surface of native code
 - Traditional memory corruption vulns
 - Similar to desktop client-side threats
- Browser/PDF/Docs = huge attack surface
- Standard Linux hardening mechanisms
 - NX, ASLR, RELRO, BIND_NOW, etc
- Android exploit mitigations are getting better...



- Android exploit mitigations have <u>slowly</u> evolved over the years...
- Before Android 2.3.x, no NX/ASLR:

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afd01000-afd02000 rw-p 00001000	1f:03 607	afd01000-afd02000 rw-p 00001000	1f:03 607
/system/lib/libstdc++.so	1f.03 487	/system/lib/libstdc++.so	1f.03 487
/system/lib/libc.so	11.05 407	/system/lib/libc.so	11.05 407
afe39000-afe3c000 rw-p 00039000	1f:03 487	afe39000-afe3c000 rw-p 00039000	1f:03 487
afe3c000-afe47000 rw-p afe3c000	00:00 0	afe3c000-afe47000 rw-p afe3c000	00:00 0
b0000000-b0013000 r-xp 00000000	1f:03 382	b0000000-b0013000 r-xp 00000000	1f:03 382
/system/bin/linker b0013000-b0014000 rw-p 00013000	1f:03 382	/system/bin/linker b0013000-b0014000 rw-p 00013000	1f:03 382
/system/bin/linker	11.05 502	/system/bin/linker	11105 502
b0014000-b001a000 rwxp b0014000	00:00 0	b0014000-b001a000 rwxp b0014000	00:00 0
bed29000-bed3e000 rwxp befeb000	00:00	be8ab000-be8c0000 rwxp befeb000	00:00 0
		#	

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- Android 2.3.x Gingerbread
 - Finally got NX support!
 - But still ineffective ASLR:





- Android 4.0 ICS
 - ASLR listed in the release notes as a new security feature!
 - But upon deeper inspection...

https://blog.duosecurity.com/2012/02/a-look-at-aslr-in-android-ice-cream-sandwich-4-0/

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- Android 4.1 Jelly Bean
 - Ok, this time we have ASLR for real...

https://blog.duosecurity.com/2012/07/exploit-mitigations-in-android-jelly-bean-4-1/

- First goal of attacker is getting a foot hold on the device with code execution
- Either by compromising an existing app or tricking user into installing a malicious app
 - Some strides made in exploit mitigations
 - More general problem of malicious apps is hard

Best to assume that malicious code/apps will be present on the user's device!

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- With code execution the attacker can:
 - Change the behavior of the exploited app
 - Steal data used by the exploited app
- The attacker can't:
 - Affected other apps on the device
 - Steal data from other apps
- Thanks to the Android "sandbox"

- Calling it a sandbox is a stretch
- Each application gets a unique uid/gid upon install

drwxr-xr-x100272048 Nov9 01:59 org.dyndns.devesh.flashlight2048 Decdrwxr-xr-x1 10046100462048 Dec8 07:18 org.freedictionary2048 Febdrwxr-xr-x1 10054100542048 Feb5 14:19 org.inodes.gus.scummvm2048 Mardrwxr-xr-x1 100392048 Mar

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- What does the privileged attack surface look like on an Android device?
 - Entire vanilla Linux kernel
 - + custom kernel modifications by Google
 - + custom drivers by third-party devs
 - Privileged system daemons (vold, etc)
 - Poorly written setuid binaries

Bottom line: Lots of attack surface to exploit!

Overview of privesc vulns

• Some vulns affect nearly all Android devices:

Name	Component	Notes
ASHMEM	kernel	custom Google mod
Exploid	init daemon	netlink source check
Gingerbreak	vold daemon	netlink source check
Levitator	kernel device driver	third-party kernel mod
Mempodroid	kernel	affected vanilla kernel
RageAgainstTheCage	adb daemon	setuid(2) return value
Wunderbar	kernel	affected vanilla kernel
ZergRush	libsysutils	memory corruption
Zimperlich	zygote	setuid(2) return value

http://jon.oberheide.org/files/bsides11-dontrootrobots.pdf

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Levitator exploit

- Levitator exploit
 - Targeted PowerVR vulnerability: /dev/pvrsrvkm
 - Allowed arbitrary kmem read/write
 - Affected popular S series devices
 - Patched in 2.3.6 after 10+ months
- Chain of custody?
 - Researcher → Google → Samsung →
 Imagination Tech → Manufacturers → Carriers

http://jon.oberheide.org/files/levitator.c

Carrier patching problem

- Carriers are terrible at patching
 - Slow, conservative patch practices
 - Inverted user/economic incentives
 - 6+ months typical vulnerability window
- One of the biggest causes of mobile insecurity
 - Carrier's tight grip on control
 - Complex ecosystem of software responsibility
 - Third-parties have no opportunity to intervene

Carrier patching problem

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Vulnerability assessment on mobile

- We can't patch the vulns*, can we at least enumerate?
- X-Ray app
 - Vulnerability assessment for Android devices
 - Launched just weeks ago
- Full stats coming next month at United Summit

http://xray.io

- If attacker escalates privileges, it's game over, right?
 - Can break out of "sandbox"
 - Tamper with applications
 - Sensitive data can be accessed/stolen
- How can we maintain security guarantees given this threat model?
 - Generally speaking, we can't!
 - But for some use cases...

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- If we're dealing with privileged attacker, assume that all system memory is compromised
 - So, key material must be kept out of memory
- HSM can provide:
 - Smartcard-style crypto engine
 - Hardware-backed tamper-proof credential storage
 - Key generation, signature computation

- SIM-based
 - mSign, etc
- SOC approahces
 - OMAP M-Shield, ARM TrustZones, etc
- Android-specific initatives
 - Google Wallet, SEEK, etc

Recent Android developments

- Keychain API released in Android 4.0
 - Primitives for credential storage

public final class KeyChain extends Object	Summary: Constants Ctors M
java.lang.Object 4 android.security.KeyChain	
Class Overview	
The KeyChain class provides access to private keys and their corre	sponding certificate chains in credential storage.
Applications accessing the $\kappa_{eychain}$ normally go through these st	eps:
1. Receive a callback from an x509KeyManager that a private key is	requested.
 Call choosePrivateKeyAlias to allow the user to select from a will be returned by the callback alias (String), or null if no priva 	list of currently available private keys and correspondin ate key is available or the user cancels the request.
 Call getPrivateKey(Context, String) and getCertificate X509KeyManager Callbacks. 	Chain(Context, String) to retrieve the credentials
An application may remember the value of a selected alias to avoid no longer valid, null will be returned on lookups using that value	prompting the user with <pre>choosePrivateKeyAlias</pre> ON S

http://nelenkov.blogspot.com.es/2011/11/using-ics-keychain-api.html

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- Further HSM support in Android 4.1
 - keymaster framework implementation
- Galaxy Nexus hardware support
 - OMAP 4, TI M-Shield platform
- Basic HSM crypto operations
 - generate_keypair, import_keypair, sign_data, verify_data, get_keypair_public, delete_keypair, delete_all

http://nelenkov.blogspot.com.es/2012/07/jelly-bean-hardware-backed-credential.html

- With a HSM, attacker can no longer extract sensitive data / key material postexploitation
- Example:
 - Duo Push private key generated on-device
 - Signatures generated within the HSM
 - Key material never leaves the credential store

Not a silver bullet: attacker can still compute signatures. Need trusted input and trusted display paths.

- Need HSM = failure of system security
 - Similar to some virtualization use cases
- Need for third-party availability
- Not the only feasible approach
 - Other avenues to contain/survive exploitation
 - On-device system containers
 - On-device hardware-level virtualization
 - Threshold crypto techniques?

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• Multiple parties are attacking the mobile security problem at a variety of layers

HSM plays an important role in device security to disrupt the attack chain even if fully compromised.

